# **Control Flow**

Chapter 6

# **Control Flow**

- Skipping most of this chapter in class
  - Skim it when reading
  - We'll just hit on a few topics that may be new
- Basic paradigms for control flow:
  - Sequencing
  - Selection
  - Iteration and Recursion
  - Procedural Abstraction
  - Concurrency
  - Nondeterminacy

### **Precedence Rules**

 Varies from language to language; the smart programmer always uses parentheses where questionable to ensure proper precedence

Fortran	Pascal	С	Ada
		++, (post-inc., dec.)	
**	not	++, (pre-inc., dec.), +, - (unary), &, * (address, contents of), !, ~ (logical, bit-wise not)	abs (absolute value) not, **
*, /	*,/, div,mod, and	<ul><li>* (binary), ∕,</li><li>% (modulo division)</li></ul>	*, /, mod, rem
+, - (unary and binary)	+, - (unary and binary), or	+, - (binary)	+, - (unary)
		<<, >> (left and right bit shift)	+, - (binary), & (concatenation)
.eq., .ne., .lt., .le., .gt., .ge. (comparisons)	<, <=, >, >=, =, <>, IN	<, <=, >, >= (inequality tests)	=, /= , <, <=, >, >=
.not.		==, != (equality tests)	
		& (bit-wise and)	
		^ (bit-wise exclusive or)	
		1 (bit-wise inclusive or)	
.and.		&& (logical and)	and, or, xor (logical operators)
.or.		[] (logical or)	
.eqv., .neqv. (logical comparisons	)	$?: (\mathrm{if} \ldots \mathrm{then} \ldots \mathrm{else})$	
		=, +=, -=, *=, /=, ½=, >>=, <<=, &=, ^=,  = (assignment)	
		. (sequencing)	

Figure 6.1: Operator precedence levels in Fortran, Pascal, C, and Ada. The operators at the top of the figure group most tightly.

**Concepts in Expression Evaluation** 

- Ordering of operand evaluation (generally none)
- · Application of arithmetic identities
  - commutativity (assumed to be safe)
  - associativity (known to be dangerous)

```
(a + b) + c works if a~=maxint and b~=minint and c<0
```

```
a + (b + c) does not
```

- Short Circuit
- · Variables as values vs. variables as references
- Orthogonality
  - Features that can be used in any combination (Algol 68)

```
a:= if b < c then d else e;  // "returns" a val
a:= begin f(b); g(c); end;  // "returns" a val
```

# **Concepts in Expression Evaluation**

- Side Effects
  - often discussed in the context of functions
  - a side effect is some permanent state change caused by execution of function
    - some noticeable effect of call other than return value
    - in a more general sense, assignment statements provide the ultimate example of side effects
      - they change the value of a variable
- Side Effects are fundamental to the von Neumann computing model
- In (pure) functional, logic, and dataflow languages, there are no such changes
  - These languages are called SINGLE-ASSIGNMENT languages
  - Expressions in a purely functional language are REFERENTIALLY TRANSPARENT

# **Concepts in Expression Evaluation**

- Boxing
  - A drawback of using the primitive value model for built-in types is they can't be passed to methods that expect references to objects
  - Boxing is "wrapping" a primitive in an object
  - Early Java: wrap explicitly using Integer, Double, etc.
  - Later Java: Automatic boxing and unboxing Integer x = 6; //6 is boxed Integer y = 2\*x + 3; //x is unboxed, 15 is boxed

# **Concepts in Expression Evaluation**

- Initialization
  - Definite assignment: variable must have a value assigned before being used
- Constructors
- Combination assignment, multiway assignment

## Structured Flow

- Sequencing
  - specifies a linear ordering on statements
    - one statement follows another
  - very imperative, Von-Neumann
- Selection
  - sequential if statements or switch/select statements

```
if ... then ... else

if ... then ... elsif ... else

(cond

(C1) (E1)

(C2) (E2)

...

(Cn) (En)

(T) (Et)

)
```

### **GOTO Statement**

- Once the primary means of control flow
  - 10 A = 1
  - 11 PRINT "HELLO"
  - 12 A = A + 1
  - 13 IF A < 100 GOTO 11
  - 14 PRINT "DONE"
- Heavily discussed in the 60's and 70's
  - Dijkstra GOTO's Considered Harmful
- Mostly abandoned in favor of structured programming; some languages allow labels (can be useful for immediate exit)
  - Not allowed at all in Java; can't even use the keyword

### Iteration

 Loops Iterators In general we may wish to iterate over the - While, repeat, for elements of any well-defined set - foreach (container or collection) Recursion ArrayList, Set (maybe), HashTable, - Tail recursion etc. - No computation follows recursive call /\* assume a, b > 0 \*/int gcd (int a, int b) { if (a == b) return a; else if (a > b) return gcd (a - b, b); else return gcd (a, b - a);}



#### Java Iterators

- Iterator used with a collection to provide sequential access to the elements in the collection
- Methods in the Iterator<T> Interface



3 public	class HashSetIteratorDemo	
4 {		
5 pt	ublic static void main(String[] args)	
6 {		
7	HashSet <string> s = new HashSet<stri< td=""><td>ng&gt;();</td></stri<></string>	ng>();
8	s.add("health");	
9	s.add("love");	
10	s.add("money");	
11	System.out.println("The set contains	:");
12	Iterator <string> i = s.iterator( );</string>	
13	<pre>while (i.hasNext( ))</pre>	
14	System.out.println(i.next( ));	
15	i.remove();	
16	System.out.println();	
17	System.out.println("The set now cont	ains:");
		- You cannot "reset" an iterator "to the
18	i = s.iterator();	beginning." To do a second iteration,
19	<pre>while (i.hasNext( ))</pre>	you create another iterator.
20	<pre>System.out.println(i.next( ));</pre>	
21	System.out.println("End of program."	');
22 }		
23 }		
Sample Dia	logue	
The set	contains:	The second second
money		The HashSet <t> object</t>
love		does not order the elements
health		it contains, but the iterator
newcell		
licatell		imposes an order on the
The set	now contains:	imposes an order on the elements.
The set money	now contains:	imposes an order on the elements.
The set money love	now contains:	imposes an order on the elements.

#### Applicative and Normal-Order Evaluation

- We normally assume all arguments are evaluated before passing them to a subroutine
  - Required for many languages so we know what to stick on the stack
  - But this need not always be the case
- Applicative Order Evaluation
  - Evaluating all arguments before the call
- Normal Order Evaluation
  - Evaluating only when the value is actually needed
  - Used in some functional languages

### Normal Order Evaluation

- Applicative Order is the norm
- Normal Order can sometimes lead to faster code, code that works when applicative-order evaluation would cause a run-time error
  - E.g. short circuit
- Lazy Evaluation in Scheme
  - Built-in functions delay and force
  - Keeps track of expressions already evaluated so it can reuse their values if they are needed more than once
    - Also called memoization
    - E.g.: fib(n) { if (n<=2) return n else return fib(n-1)+fib(n-2) }