# Functional Programming in Scheme 

## Functional Programming

- Online textbook: http://www.htdp.org/
- Original functional language is LISP
- LISt Processing
- The list is the fundamental data structure
- Developed by John McCarthy in the 60's
- Used for symbolic data processing
- Example apps: symbolic calculations in integral and differential calculus, circuit design, logic, game playing, AI
- As we will see the syntax for the language is extremely simple
- Scheme
- Descendant of LISP


## Functional Languages

- "Pure" functional language
- Computation viewed as a mathematical function mapping inputs to outputs
- No notion of state, so no need for assignment statements (side effects)
- Iteration accomplished through recursion
- In practicality
- LISP, Scheme, other functional languages also support iteration, assignment, etc.
- We will cover some of these "impure" elements but emphasize the functional portion
- Equivalence
- Functional languages equivalent to imperative
- Core subset of C can be implemented fairly straightforwardly in Scheme
- Scheme itself implemented in C
- Church-Turing Thesis


## Lambda Calculus

- Foundation of functional programming
- Developed by Alonzo Church, 1941
- A lambda expression defines
- Function parameters
- Body
- Does NOT define a name; lambda is the nameless function. Below $x$ defines a parameter for the unnamed function:

$$
(\lambda x \cdot x * x)
$$

## Lambda Calculus

- Given a lambda expression

$$
(\lambda x \cdot x * x)
$$

- Application of lambda expression

$$
((\lambda x \cdot x * x) 2) \rightarrow 4
$$

- Identity
$(\lambda x \cdot x)$
- Constant 2: $(\lambda x \cdot 2)$


## Lambda Calculus

- Any identifier is a lambda expression
- If M and N are lambda expressions, then the application of $M$ to $N,(M N)$ is a lambda expression
- An abstraction, written $(\lambda x \cdot M)$ where $x$ is an identifier and M is a lambda expression, is also a lambda expression


## Lambda Calculus

$$
\begin{aligned}
\text { LambdaExpression } & \rightarrow \text { ident }|(M N)|(\lambda \text { ident } \cdot M) \\
M & \rightarrow \text { LambdaExpression } \\
N & \rightarrow \text { LambdaExpression }
\end{aligned}
$$

Examples

$$
\begin{aligned}
& x \\
& (\lambda x \cdot x) \\
& ((\lambda x \cdot x)(\lambda y \cdot y))
\end{aligned}
$$

## Lambda Calculus

First Class Citizens

- Functions are first class citizens
- Can be returned as a value
- Can be passed as an argument
- Can be put into a data structure as a value
- Can be the value of an expression

$$
\begin{aligned}
& ((\lambda x \cdot x * x)(\lambda y \cdot 2))=(\lambda x \cdot 2 * 2)=4 \\
& ((\lambda \mathrm{x} \cdot(\lambda \mathrm{y} \cdot \mathrm{x}+\mathrm{y})) 21)=((\lambda y \cdot 2+\mathrm{y}) 1)=3
\end{aligned}
$$

## Lambda Calculus

## Functional programming is essentially an applied lambda calculus with built in <br> - constant values <br> - functions

E.g. in Scheme, we have ( $* x$ x ) for $x^{*} x$ instead of $\lambda x \cdot x^{*} x$

## Functional Languages

- Two ways to evaluate expressions
- Eager Evaluation or Call by Value
- Evaluate all expressions ahead of time
- Irrespective of if it is needed or not
- May cause some runtime errors
- Example
(foo 1 (/ 1 x ))
Problem; divide by 0


## Lambda Calculus

- Lazy Evaluation
- Evaluate all expressions only if needed
(foo $1(/ 1 \mathrm{x})$ ) ; (/ 1 x ) not needed, so never eval'd
- Some evaluations may be duplicated
- Equivalent to call-by-name
- Allows some types of computations not possible in eager evaluation
- Example
- Infinite lists
- E.g.. Infinite stream of 1 's, integers, even numbers, etc.
- Replaces tail recursion with lazy evaluation call
- Possible in Scheme using (force/delay)


## Running Scheme for Class

- A version of Scheme called Racket (formerly PLT/Dr Scheme) is available on the Windows machines in the NSB/ENGR Labs
- Download: http://racket-lang.org/
- Unix, Mac versions also available if desired


## Racket

- You can type code directly into the interpreter and Scheme will return with the results:



## Make sure right Language is selected



Use the "Pretty<br>Big" language<br>choice - it is<br>closer to Scheme<br>than others

Welcome to DrRacket, version 5.2.1[3m]
anguage- Beginning Student; memory limit: 128 MB
$>(1$ ambda (x) (+1 x) 1)
lambda: found a lambda that is not a function definition

## $>$

## Racket - Loading Code

- You can open code saved in a file. Racket uses the extension ".rkt" so consider the following file "factorial.rkt" created with a text editor or saved from Racket:
factorial (lambda (n)
(cond
((=n 1) 1)
(else (* n (factorial (- n 1))))
)


## Functional Programming Overview

- Pure functional programming
- No implicit notion of state
- No need for assignment statement
- No side effect
- Looping
- No state variable
- Use Recursion
- Most functional programming languages have side effects, including Scheme
- Assignments
- Input/Output


## Scheme Programming Overview

- Refreshingly simple
- Syntax is learned in about 10 seconds
- Surprisingly powerful
- Recursion
- Functions as first class objects (can be value of an expression, passed as an argument, put in a data structure)
- Implicit storage management (garbage collection)
- Lexical scoping
- Earlier LISPs did not do that (dynamic)
- Interpreter
- Compiled versions available too


## Expressions

- Syntax - Cambridge Prefix
- Parenthesized
-(* 34 )
$-(*)+23) 5)$
-(f 34 )
- In general:
-(functionName arg1 arg2 ...)
- Everything is an expression
- Sometimes called s-expr (symbolic expr)


## Expression Evaluation

- Replace symbols with their bindings
- Constants evaluate to themselves
- 2, 44, \#f
- No nil in Racket; use '()
- Nil = empty list, but Racket does have empty
- Lists are evaluated as function calls written in Cambridge Prefix notation
(+23)
(* +2 3) 5)


## Scheme Basics

## - Atom

- Anything that can't be decomposed further
- a string of characters beginning with a letter, number or special character other than (or )
- e.g. 2, \#t, \#f, "hello", foo, bar
- \#t = true
- \#f = false
- List
- A list of atoms or expressions enclosed in ()
- (), empty, (1 23 ), (x (2 3)), (()()())


## Scheme Basics

- S-expressions
- Atom or list
- () or empty
- Both atom and a list
- Length of a list
- Number at the top level


## Quote

- If we want to represent the literal list (abc)
- Scheme will interpret this as apply the arguments $b$ and $c$ to function $a$
- To represent the literal list use "quote"
$-($ quote x$) \rightarrow \mathrm{x}$
$-($ quote $(\mathrm{abc})) \rightarrow(\mathrm{abc})$
- Shorthand: single quotation mark
' $\mathrm{a}==$ (quote a )
$'(a b c)=($ quote $(a b c))$


## Global Definitions

- Use define function
(define f 20)
(define evens '(0 246 8))
(define odds '(13579))
(define color 'red)
(define color blue) ; Error, blue undefined
(define num f) ; num $=20$
(define num 'f) ; symbol f
(define s "hello world") ; String


## Lambda functions

- Anonymous functions
- (lambda (<formals>) <expression>)
- (lambda (x) (* x x))
$-((\operatorname{lambda}(x)(* x x)) 5) \rightarrow 25$
- Motivation
- Can create functions as needed
- Temporary functions : don't have to have names
- Can not use recursion


## Named Functions

- Use define to bind a name to a lambda expression
(define square (lambda (x) (* x x)))
(square 5)
- Using lambda all the time gets tedious; alternate syntax:
(define (<function name> <formals>) <expression1> <expression2> ...)
Last expression evaluated is the one returned
(define (square x ) (* x x ) )
(square 5) $\rightarrow 25$


## Conditionals

```
(if <predicate> <expression1> <expresion2>)
- Return value is either expr1 or expr2
(cond (P1 E1)
(P2 E2)
\(\left(\mathrm{P}_{\mathrm{n}} \mathrm{E}_{\mathrm{n}}\right)\)
(else \(\mathrm{E}_{\mathrm{n}+1}\) ))
```

- Returns whichever expression is evaluated


## Common Predicates

- Names of predicates end with ?
- Number? : checks if the argument is a number
- Symbol? : checks if the argument is a symbol
- Equal? : checks if the arguments are structurally equal
- Null? : checks if the argument is empty
- Atom? : checks if the argument is an atom
- Appears undefined in Racket but can define ourselves
- List? : checks if the argument is a list


## Conditional Examples

- (if (equal? 12) 'x ‘y) ; y
- (if (equal? 2 2) 'x'y) ; x
- (if (null? '()) 122 ; 1
- (cond
((equal? 12 2) 1)
((equal? 23 ) 2)
(else 3)) ; 3
- (cond
((number? 'x) 1)
((null? 'x) 2)
$(($ list? ' $(\mathrm{ab} \mathrm{c}))(+23)) \quad ; 5$
)


## Dissecting a List

- Car : returns the first argument
- (car '(2 34 ) )
- (car '((2) 44$)$ )
- Defined only for non-null lists
- Cdr : (pronounced "could-er") returns the rest of the list
- Racket: list must have at least one element
- Always returns a list
- (cdr '(2 3 4))
- (cdr ‘(3))
- ( cdr ' ${ }^{(((3)))))}$
- Compose
- (car (cdr ‘(4 5 5)))
- (cdr (car $\left.\left.{ }^{〔}((34))\right)\right)$


## Shorthand

- $(\operatorname{cadr} x)=(\operatorname{car}(\operatorname{cdr} x))$
- $(\operatorname{cdar} \mathrm{x})=(\operatorname{cdr}(\operatorname{car} \mathrm{x}))$
- $(\operatorname{caar} \mathrm{x})=(\operatorname{car}(\operatorname{car} \mathrm{x}))$
- $(\operatorname{cddr} x)=(c d r(c d r x))$
- $(\operatorname{cadar} \mathrm{x})=(\operatorname{car}(\operatorname{cdr}(\operatorname{car} \mathrm{x})))$
- ... etc... up to 4 levels deep in Racket
- $($ cddadr $x)=$ ?


## Why Car and Cdr?

- Leftover notation from original implementation of Lisp on an IBM 704
- CAR = Contents of Address part of Register
- Pointed to the first thing in the current list
- $\mathrm{CDR}=$ Contents of Decrement part of Register
- Pointed to the rest of the list


## Building a list

## - Cons

- Cons(truct) a new list from first and rest
- Takes two arguments
- Second should be a list
- If it is not, the result is a "dotted pair" which is typically considered a malformed list
- First may or may not be a list
- Result is always a list


## Building a list

$X=2$ and $Y=(345):($ cons $x y) \rightarrow$
(2345)
$\mathrm{X}=()$ and $\mathrm{Y}=(\mathrm{abc}):($ cons $\mathrm{x} y) \rightarrow$
(() abc)
$\mathrm{X}=\mathrm{a}$ and $\mathrm{Y}=():(\operatorname{cons} \mathrm{x} y) \rightarrow$
(a)

- What is
- (cons 'a (cons 'b (cons 'c '())))
- (cons (cons 'a (cons 'b ‘())) (cons 'c ‘()))


## Numbers

- Regular arithmetic operators are available
$+,-, *, /$
- May take variable arguments (+234), (* 459 11)
- (/ 9 2) $\rightarrow 4.5$; (quotient 92$) \rightarrow 4$
- Regular comparison operators are available $<><=>=$
- E.g. $(=5(+32)) \quad \rightarrow$ \#t
$=$ only works on numbers, otherwise use equal?


## Example

- Sum all numbers in a list

```
(define (sumall list)
    (cond
    ((null? list) 0)
    (else (+ (car list) (sumall (cdr list))))))
```

Sample invocation: (sumall '(3 45 1))

## Example

## - Make a list of n identical values

```
(define (makelist n value)
    (cond
        ((= n 0)'())
        (else
                (cons value (makelist (- n 1) value))
    )
)
)
```

In longer programs, careful matching parenthesis.

## Example

- Determining if an item is a member of a list

```
(define (member? item list)
    (cond ((null? list) #f)
            ((equal? (car list) item) #t)
            (else (member? item (cdr list)))
        )
    )
```

Scheme already has a built-in (member item list) function that returns the list after a match is found

## Example

## - Remove duplicates from a list

```
(define (remove-duplicates list)
    (cond ((null? list) '())
        ((member? (car list) (cdr list))
        (remove-duplicates (cdr list)))
        (else
        (cons (car list) (remove-duplicates (cdr list))))
    )
)
```

